Exam Automated Reasoning

Thursday, 3 February 2005, 9 - 12 h.

NB. The exam will be corrected and graded by humans, not by a computer. Therefore, you need not to bother too much about the syntactical peculiarities of PVS and Promela.

1. Let R be a binary symmetrical relation over some set X, so $R \subseteq X^2$ and $\forall x, y \in X((x, y) \in R \to (y, x) \in R)$. Then the following property holds:

$$(R^{-1})^* \subseteq R^*$$

Here $^{-1}$ denotes the inverse, so $R^{-1} = \{(y, x) \mid (x, y) \in R\}$, and * denotes transitive closure:

$$R^* = \{ (f(0), f(n)) \mid n \in N \land f : N \to X \land \forall i (i \le n \to (f(i), f(i+1)) \in R) \},\$$

Formulate a PVS-theory and a lemma that expresses this property. (You may use predicates to model relations.) Then sketch how this lemma can be proved with PVS.

2. Consider three processes, defined by

```
process Thread (self := 0 to 2)
var priv: int := 0
do :: true ->
     RW
     priv ++
     synch.self
od
```

The idea is that the processes do relevant work in RW, but need to be synchronized in such a way that

 $\begin{array}{l} priv.0 \leq priv.1+1 \ , \\ priv.0 + priv.1 \leq priv.2+2 \ , \\ priv.2 \leq priv.0 + priv.1+1 \ . \end{array}$

always hold.

- (a) Implement the three commands sych.0, sych.1, sych.2 to realize this, under the following conditions:
 - i. unnecessary waiting is to be avoided;
 - ii. the commands are not allowed to modify priv;
 - iii. they may only contain assignments, atomic waits and busy-waiting loops;
 - iv. you may only use atomic commands that are *simply shared*, i.e. they may refer to at most one shared variable at most once.
- (b) Show that your solution admits an execution where the value of *priv.2* becomes arbitrary large.
- (c) Indicate how you may check the correctness of your solution with Spin. Pay attention to correctness, deadlock and progress. Do not forget to reduce your solution to a finite state space.

Hint: introduce auxiliary variables, e.g. shared variables sh and a private variable own.